

Service Level Agreement Control in the Presence of Heterogeneous Traffic and QoS Requirements

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Abstract – The paper proposes a novel measurement-based *Equivalent Bandwidth* technique that computes the bandwidth to be allocated to a buffer which conveys heterogeneous traffic (both concerning traffic sources and QoS requirements), without using any closed-form expression. The effectiveness of the algorithm is checked through simulation analysis.

Keywords – QoS, Equivalent bandwidth, Measurement control.

I. INTRODUCTION

Equivalent bandwidth (EqB) is defined as the minimum service rate to be provided to a traffic buffer to guarantee a certain degree of *Quality of Service* (QoS) in terms of objective parameters (packet loss, delay, jitter). EqB techniques are obtained analytically for homogeneous traffic trunks, with respect to a single QoS constraint. Modern network solutions often imply the aggregation of service classes with different QoS constraints, thus generating heterogeneous trunks from the point of view of both traffic sources and QoS requirements. This situation leads to the need to develop new equivalent bandwidth techniques so as to match heterogeneity.

II. ASSUMPTIONS

There are N traffic classes. $a_i(t)$ is the *input rate* process of the i -th traffic class and $a(t)$ the aggregate process of all $a_i(t)$, $i = 1, \dots, N$. Traffic is conveyed towards a single buffer, modelled through a *Stochastic Fluid Model* [1, 2]. $a(t)$ is supposed ergodic for now, so that a single realization is representative of the entire process. This assumption will be relaxed later. There is no knowledge of $a_i(t)$ processes, as well as of the aggregate process $a(t)$. Additionally, aggregation may involve also buffering and encapsulation operations as typically done in real network nodes. It makes $a(t)$ analytical modelling virtually impossible to get also in case of full knowledge of $a_i(t)$ processes. The only information about $a_i(t)$ and $a(t)$ may be got through real measures. The *service rate* of the buffer is $R(t)$. $l_i(R(t), t)$ is the *loss rate* process of the i -th traffic class, measured in [bps]. The average value of the loss rate is defined in (1) for $i = 1, \dots, N$. No analytical expression for $l_i(R(t), t)$ is supposed available. Information about its behaviour is got by measures. The entire system model is reported in Fig. 1.

The SLA (*Service Level Agreement*) for each traffic class is composed of a *Packet Loss Probability* threshold (PLP_i^*). It means that the amount of *feasible loss rate* must be limited in

any time instant by the process $l_i^*(t) = PLP_i^* \cdot a_i(t)$, measured in [bps], whose average value is contained in (2) for $i = 1, \dots, N$.

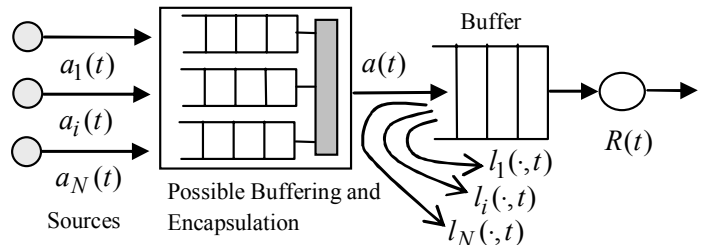


Fig. 1. System model.

$$\bar{l}_i = \lim_{\tau \rightarrow \infty} \frac{1}{\tau} \int_{\tau} l_i(t) dt \quad (1) \quad \bar{l}_i^* = \lim_{\tau \rightarrow \infty} \frac{1}{\tau} \int_{\tau} l_i^*(t) dt \quad (2)$$

III. PROBLEM DEFINITION AND SOLUTION

The aim is to provide the minimum buffer service rate so that the maximum quadratic distance between \bar{l}_i and \bar{l}_i^* is minimized. It corresponds to define the optimization problem in (3), introduced in this paper and identified as *Generalized Equivalent Bandwidth* (GEqB).

$$R^* = \arg \min_R l_{\Delta}(\cdot, R), \quad l_{\Delta}(\cdot, R) = \text{Max}_i [\bar{l}_i - \bar{l}_i^*]^2 \quad (3)$$

Being the involved stochastic processes unknown, GEqB problem is solved by taking measures over a given k -th *observation horizon* (OH), $T_k = [t_{k-1}, t_k]$, $k = 1, 2, \dots$, and performing a sequence of bandwidth reallocations $R(t_k)$, $k = 1, 2, \dots$, each T_k , based on the gradient method. The loss rate $l_i(R(t_k), t)$ and the feasible loss rate $l_i^*(t)$ processes are averaged over each OH, giving origin to the quantities \hat{l}_i^k in (4) and $\hat{l}_i^{*,k}$ in (5). Being used to solve the GEqB problem, \hat{l}_i^k and $\hat{l}_i^{*,k}$ must be representative of the average values \bar{l}_i and \bar{l}_i^* , $\forall i = 1, \dots, N$ and $\forall k$.

$$\hat{l}_i^k = \frac{1}{T_k} \int_{T_k} l_i(t) dt \quad (4) \quad \hat{l}_i^{*,k} = \frac{1}{T_k} \int_{T_k} l_i^*(t) dt \quad (5)$$

Bandwidth allocation at instant T_k is ruled by the algorithm introduced in Fig. 2 and called *Gradient-based Generalized Equivalent Bandwidth* (G^2EqB) algorithm. $step_k$ is the gradient stepsize. Condition a) means that the allocated bandwidth needs to be increased (i.e., $\hat{l}_i^{(k+1)} - \hat{l}_i^{*,(k+1)}$ for some $i = 1, \dots, N$). The $Max[\cdot]$ operator in **a**) allows exploiting the largest bandwidth need deriving from the currently most “suffering” traffic class. The state of surffrance is measured by cost sensitivity $2 \cdot \frac{\partial \hat{l}_i(R)}{\partial R} \Big|_{R=R(t_k)} [\hat{l}_i^{*,(k+1)} - \hat{l}_i^{(k+1)}]$.

Condition b) states the opposite; the $min[\cdot]$ operator leads to a bandwidth reduction with respect to the traffic class whose sensitivity distance from the performance constraint (i.e., $2 \cdot \frac{\partial \hat{l}_i(R)}{\partial R} \Big|_{R=R(t_k)} [\hat{l}_i^{(k+1)} - \hat{l}_i^{*,(k+1)}]$) is the smallest one. In other words, when bandwidth is increased (condition a)), largest bandwidth steps are followed; when bandwidth is decreased (condition b)), smallest steps are used.

Derivatives $\frac{\partial \hat{l}_i(R)}{\partial R}$ represent the sensitivity of the loss to infinitesimal variations of the rate serving the buffer. Intuitively they depend on the speed with which the system passes from an empty to a full state. They can be obtained by observing the buffer state evolution within each OH, which is divided into N_{T_k} busy periods (where the buffer is not empty) identified by the variable bp . If there is one traffic class, the derivative exact form is presented in [1] but it is still unknown in the multiple class case when the service rate is the control variable. This paper introduces the approximation in (6). (6) is equality in case of single class, as proved in [1]. $[{}^i at_{T_k}^{bp}(R(t_{k-1})) - {}^i ll_{T_k}^{bp}(R(t_{k-1}))]$ is the contribution to information loss of the i -th traffic class for the busy period bp within T_k , $k = 1, 2, \dots$. ${}^i at_{T_k}^{bp}$ is the arrival time of the first packet of service class i within the busy period bp . ${}^i ll_{T_k}^{bp}$ is the time when the last loss of class i occurs during bp .

$$\frac{\partial \hat{l}_i(R)}{\partial R} \Big|_{R=R(t_{k-1})} \cong -\frac{1}{T_k} \sum_{bp=1}^{N_{T_k}} [{}^i at_{T_k}^{bp}(R(t_{k-1})) - {}^i ll_{T_k}^{bp}(R(t_{k-1}))] \quad (6)$$

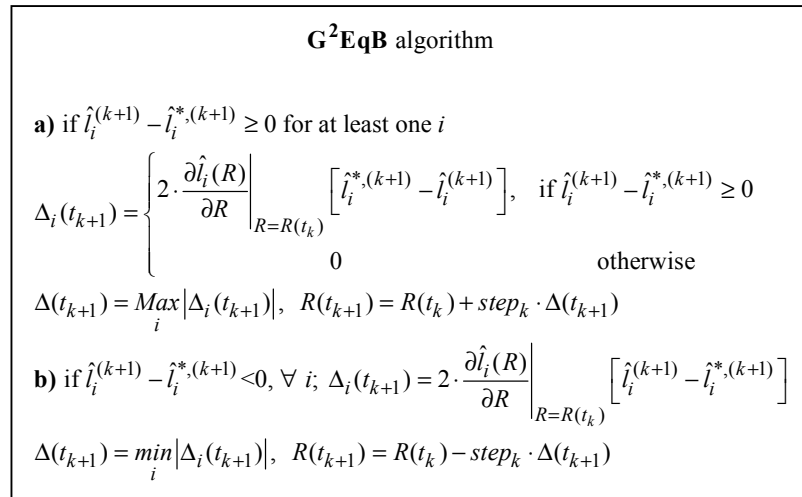


Fig. 2. G^2EqB algorithm.

IV. ALGORITHM CONVERGENCE

Technical conditions for convergence to global optimum ($R(t_k) \xrightarrow{k \rightarrow \infty} R^*$) are: **1)** Ergodic stochastic processes, **2)** decreasing behaviour of $step_k$, **3)** gradient bounded within the control domain $R(t) \in \mathfrak{R}^+$, $\forall t$, and **4)** non-existence of local optima. **1)** and **2)** are assumptions. Concerning **3)**, the lengths of buffer busy periods are bounded by OH size; measured loss rate at the end of each OH cannot be infinite. Concerning **4)**, the loss rate of a traffic queue can be reasonably assumed to be continuous, differentiable, with a negative derivative with respect to the service rate, so the cost function is also continuous, differentiable with unique minimum. The convergence speed depends on the length of OH and on the gradient stepsize. The length of OH is

important because, on one hand, it must be long enough to assure that \hat{l}_i^k and $\hat{l}_i^{*,k}$ are representative of the average values \bar{l}_i and \bar{l}_i^* , $\forall i = 1, \dots, N$ and $\forall k$, but, on the other hand, it must be short enough to assure quick convergence. In this context, the assumption of process ergodicity may be relaxed and limited to the time that the sequence of bandwidth allocations $R(t_k)$, $k = 1, 2, \dots$ needs to converge to R^* . When $a(t)$ changes its statistical behaviour, a new $GEqB$ problem is solved by supposing $a(t)$ ergodic at least within convergence time and by starting G^2EqB algorithm again. In consequence, tuning OH length is important also to get fast reactions to traffic variations. Concerning gradient stepsize dimension, tests not reported here have shown that it is not a critical parameter for convergence. Actually, convergence condition

2) may be relaxed: setting a proper constant value of the gradient stepsize as done in [2] and in this paper, does not affect convergence. For example if $step_k = 1$ in the tests reported below the algorithm converges but stepsize length adaptation helps improve convergence speed and limit bandwidth oscillations.

V. PERFORMANCE ANALYSIS AND DISCUSSION

Due to the complexity of the overall input rate process $a(t)$, equivalent bandwidth approaches which use complex mathematical descriptors may be hardly applied in real time. The approach in [3] (called EqB in the following) is applicable in this context and used here to make a comparison with G^2EqB . $m_a(t_k)$ and $\sigma_a(t_k)$ are the measured mean and standard deviation of $a(t)$ over the k -th OH. Bandwidth is assigned at time $t_k, k = 1, 2, \dots$ as in (7). PLP_{EqB}^* is the allowed PLP upper bound and is defined as the most stringent PLP requirement out of N SLAs.

$$R(t_k) = m_a(t_k) + d \cdot \sigma_a(t_k), \quad d = \sqrt{-2 \ln(PLP_{EqB}^*) - \ln(2\pi)} \quad (7)$$

A. G^2EqB versus EqB: rate provision and convergence

VoIP SLA is considered. Each source is an on-off process. Mean on and off time durations are exponentially distributed with mean 1.008 s and 1.587 s, respectively. Peak bandwidth is 16 kbps. VoIP traffic enters an IP buffer whose length and service rate (set by the traffic peak bandwidth) guarantee no packet loss rate. IP traffic is encapsulated in ATM (via AAL5) so generating the process $a(t)$ as output of the “Buffering and Encapsulation” box in Fig. 1. $a(t)$ enters the ATM buffer (1600 bytes), where the VoIP loss rate \hat{l}^k in IP packets is measured each T_k . PLP_{VoIP}^* is set to $2 \cdot 10^{-2}$; G^2EqB OH to

30 s; gradient stepsize to 6.0. EqB OH is either fixed to 30 s or tuned through an approximation of the Dominant Time Scale principle that computes the optimal OH size to get a reliable estimate of EqB statistics. The two alternatives are respectively identified as “EqB OH 30s” and “EqB OH DTS”, in the following. Fig. 3 and Fig. 4 show, respectively, PLP and corresponding allocated bandwidth of G^2EqB and EqB. The number of VoIP sources is increased of 10 from 70 to 110 each 3000 seconds. Average PLP results are: $4.40 \cdot 10^{-3}$ for G^2EqB and $1.18 \cdot 10^{-2}$ for EqB OH 30s. EqB OH DTS assures null packet loss. Average allocated bandwidths are: 0.842 Mbps for G^2EqB , 0.867 Mbps for EqB OH 30s, and 1.31 Mbps for EqB OH DTS. Even if the average PLP values seem to be satisfying for all the schemes, the simple observation of Figs. 3 and 4 suggests that G^2EqB reacts quickly to traffic changes also minimizing bandwidth oscillations. EqB OH DTS always matches PLP_{VoIP}^* request but implies a relevant bandwidth waste; EqB OH 30s often fails to guarantee PLP_{VoIP}^* and introduces wide bandwidth and PLP oscillations. Quantitative metrics may help the interpretation of this qualitative behaviour. PLP standard deviation is $7.4 \cdot 10^{-3}$ for G^2EqB and $1.33 \cdot 10^{-2}$ for EqB OH 30s. The percentage of the OH periods where PLP is over threshold is 5% for G^2EqB and 18.6% for EqB OH 30s. The average difference value between measured PLP and PLP_{VoIP}^* selecting the OH periods where PLP is over threshold is $4.22 \cdot 10^{-4}$ for G^2EqB and $2.77 \cdot 10^{-3}$ for EqB OH 30s. G^2EqB allows minimizing the distance between measured and threshold values.

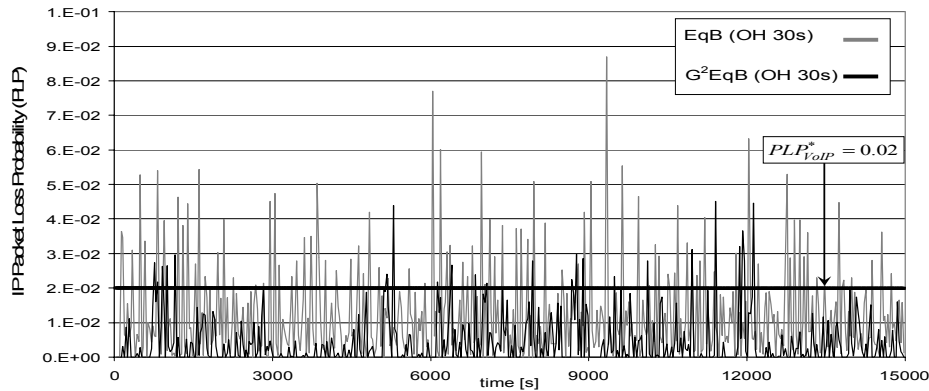


Fig. 3. G^2EqB and EqB: PLP.

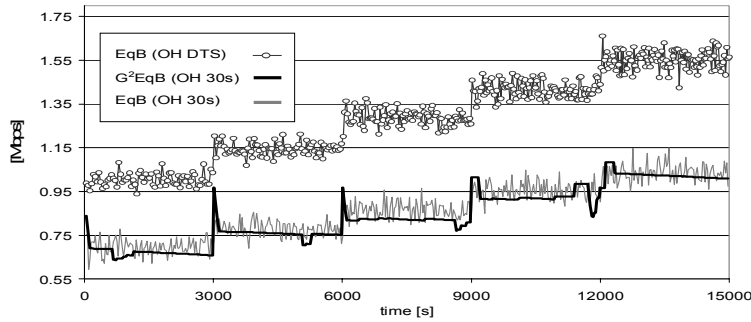


Fig. 4. G2EqB versus EqB: bandwidth allocation.

B. G²EqB versus EqB: heterogeneous traffic

A real video trace (“Jurassik park” from [4]) is added to the VoIP scenario used in section V.A. Peak and average rate are 1.418 and 0.280 Mbps. Video enters an IP buffer whose length and service rate (set to 1.418 Mbps) guarantee no loss. Both VoIP and video traffic are encapsulated over DVB (packets of 188 bytes) at the exit of the IP buffers and generate the process $a(t)$. PLP_{video}^* is set to $5 \cdot 10^{-3}$. Table 1 contains the average measured video and VoIP PLPs together with the allocated bandwidth to the DVB buffer in [Mbps] for G²EqB OH 3min (G²EqB in Table 1), EqB OH 3min, and

EqB OH DTS. DVB buffer dimension is changed as well as the number of VoIP calls. Each single test simulates 107 overall minutes. The average G²EqB PLP is always close to but below the threshold of the most restrictive requirement (video PLP). G²EqB is adaptive to buffer length because its behaviour depends only on loss measures. EqB-based schemes do not adapt to buffer length: EqB OH 3min underestimates the bandwidth and fails to match video requirements. EqB OH DTS behaves similarly for short buffer length while overestimates the bandwidth for larger buffer dimensions.

Number VoIP Calls	Buffer Length [bytes]	Allocated Bandwidth G ² EqB / EqB OH 3min / EqB OH DTS	Video PLP G ² EqB / EqB OH 3min / EqB OH DTS	VoIP PLP G ² EqB / EqB OH 3min / EqB OH DTS
30	9400	0.94 / 0.79 / 0.87	$2.62 \cdot 10^{-3}$ / $8.26 \cdot 10^{-2}$ / $1.05 \cdot 10^{-2}$	$1.20 \cdot 10^{-4}$ / $5.58 \cdot 10^{-3}$ / $5.93 \cdot 10^{-4}$
30	18800	0.84 / 0.79 / 0.87	$3.20 \cdot 10^{-3}$ / $2.96 \cdot 10^{-2}$ / 0.0	$2.22 \cdot 10^{-4}$ / $2.13 \cdot 10^{-3}$ / 0.0
30	28200	0.82 / 0.79 / 0.87	$1.78 \cdot 10^{-3}$ / $1.34 \cdot 10^{-2}$ / 0.0	$1.24 \cdot 10^{-4}$ / $9.14 \cdot 10^{-4}$ / 0.0
60	9400	1.47 / 1.26 / 1.41	$3.94 \cdot 10^{-3}$ / $8.26 \cdot 10^{-2}$ / $1.15 \cdot 10^{-2}$	$2.26 \cdot 10^{-4}$ / $7.77 \cdot 10^{-3}$ / $5.63 \cdot 10^{-4}$
60	18800	1.41 / 1.26 / 1.41	$2.23 \cdot 10^{-3}$ / $5.13 \cdot 10^{-2}$ / $1.54 \cdot 10^{-3}$	$1.55 \cdot 10^{-4}$ / $3.96 \cdot 10^{-3}$ / $8.17 \cdot 10^{-5}$
60	28200	1.39 / 1.26 / 1.41	$1.38 \cdot 10^{-3}$ / $3.0 \cdot 10^{-2}$ / $4.20 \cdot 10^{-4}$	$1.04 \cdot 10^{-4}$ / $2.32 \cdot 10^{-3}$ / $8.17 \cdot 10^{-5}$
90	9400	1.97 / 1.73 / 2.00	$4.13 \cdot 10^{-3}$ / $1.07 \cdot 10^{-1}$ / $1.58 \cdot 10^{-2}$	$2.38 \cdot 10^{-4}$ / $8.11 \cdot 10^{-3}$ / $8.28 \cdot 10^{-4}$
90	18800	1.87 / 1.73 / 2.0	$3.26 \cdot 10^{-3}$ / $5.33 \cdot 10^{-2}$ / $4.35 \cdot 10^{-3}$	$1.94 \cdot 10^{-4}$ / $3.97 \cdot 10^{-3}$ / $2.37 \cdot 10^{-4}$
90	28200	1.84 / 1.73 / 2.0	$1.73 \cdot 10^{-3}$ / $3.07 \cdot 10^{-2}$ / $2.02 \cdot 10^{-3}$	$1.02 \cdot 10^{-4}$ / $2.35 \cdot 10^{-3}$ / $2.37 \cdot 10^{-4}$

Table. 1. G²EqB and EqB: packet loss and bandwidth allocation.

VI. CONCLUSIONS AND FUTURE WORK

A novel equivalent bandwidth algorithm to automatically adapt the rate assigned to a buffer and counteract time varying system conditions is introduced. It is based only on measures. No closed-form expressions, no a-priori information about source statistical properties, and no assumptions about buffer dimension are imposed. To achieve quick on-line convergence, the algorithm requires: proper off line dimensioning of gradient descent stepsize and proper length of the observation horizon. Time-scale of traffic changes must be slower than the time required for convergence. Otherwise, the resulting allocations are suboptimal.

Future research may regard: **1)** multi-objective control, namely, joint control of non homogeneous (sometimes conflicting) performance metrics (loss versus delay, versus delay jitter of the packets); **2)** implementation and validation of the algorithm within a Linux-based testbed architecture [5]; **3)** tuning of algorithm parameters to reliably support QoS in heterogeneous conditions; **4)** performance analysis of the control algorithm with respect to different traffic categories, for example in the presence of congestion control (e.g., TCP).

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